

Standard Form Polyhedra

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Polyhedra in standard form

generic polyhedron

$$\mathcal{P} = \left\{ x \mid \begin{array}{l} Ax = b \\ Cx \leq d \end{array} \right\}$$

standard-form polyhedron

$$\mathcal{P} = \left\{ x \mid \begin{array}{l} Ax = b \\ x \geq 0 \end{array} \right\} \quad \text{with} \quad b \geq 0$$

Converting to standard form: positive b

- elements from both b and d will b in standard form.

For $b_i < 0$, replace

$$a_i x = b_i \longrightarrow (-a_i)x = (-b_i)$$

For $d_i < 0$, replace

$$c_i x \leq d_i \longrightarrow (-c_i)x \geq (-d_i)$$

$$c_i x \geq d_i \longrightarrow (-c_i)x \leq (-d_i).$$

Converting to standard form: free variables

- x_i is called a **free variable** if it has no constraints
- there are no free variables in standard form — every variable must be nonnegative

Converting free variables

- every free variable x_i is replaced with two new variables x'_i and x''_i , ie,

$$x_i := x'_i - x''_i, \quad x'_i \geq 0 \quad \text{and} \quad x''_i \geq 0$$

- x'_i encodes the positive part of x_i
- x''_i encodes the negative part of x_i

Converting to standard form: slack and surplus

For every inequality constraint of the form

$$c_i x \leq d_i \quad (c_i x \geq d_i)$$

introduce a new **slack** (or **surplus**) variable s_i , replacing the inequality with two constraints

$$\begin{array}{ll} c_i x + s_i = d_i & \left(\begin{array}{l} c_i x - s_i = d_i \\ s_i \geq 0 \end{array} \right) \\ s_i \geq 0 & \end{array}$$

Basic solutions in standard form

x^* is a **basic solution** if the vectors

$$a_{i_1}, a_{i_2}, \dots, a_{i_n}, \quad i_j \in \mathcal{B}$$

are linearly independent

In standard form, there are

- n variables (x_1, \dots, x_n)
- $m + n$ total constraints
 - m equality constraints $(Ax = b)$
 - n inequality constraints $(x \geq 0)$

for any basic solution x ,

- the basic set \mathcal{B} must have n elements
- thus, exactly n of the constraints need to be active at x
- m equality constraints are always satisfied
- thus $n - m$ of the inequality constraints $x \geq 0$ should be “active”

Basic solutions in standard form

Choosing $n - m$ of the inequality constraints to be active is the same as choosing $n - m$ variables x_i to be zero. Making x_i zero effectively eliminates column i from the matrix A .

This is equivalent to choosing m columns of A ! To be a basic solution, we also need these m columns to be linearly independent. So, permute the variables and partition

$$AP = [B \quad N] \quad \text{where } B \text{ is nonsingular}$$

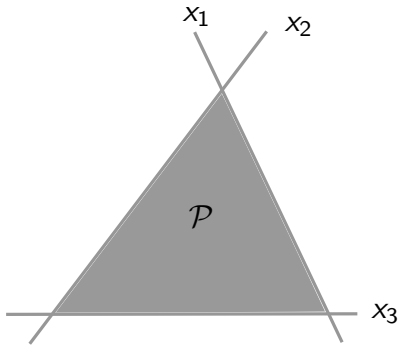
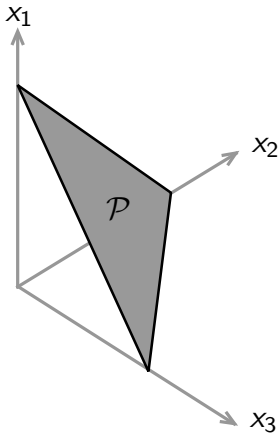
Now we have

$$\bar{A}x = \begin{bmatrix} B & N \\ I \end{bmatrix} \begin{bmatrix} x_B \\ x_N \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix}$$

$$x_N = 0$$

$$Bx_B = b$$

Two-dimensional representation



Degeneracy: inequality form

polyhedron in inequality form:

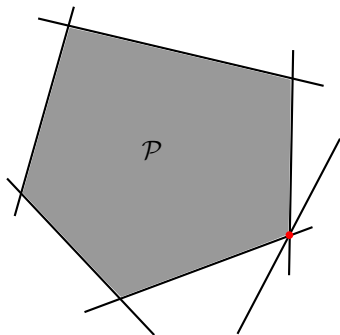
$$Ax \leq b$$

a basic feasible solution x^* with

$$a_i^T x^* = b_i, \quad i \in \mathcal{B} \quad \text{and} \quad a_i^T x^* < b_i, \quad i \notin \mathcal{B}$$

is **degenerate** if $\#$ of indices in \mathcal{B} is greater than n

- property of the **description** of the polyhedron
- affects the performance of some algorithms
- disappears for small perturbations of b



Degeneracy: standard form

polyhedron in standard form:

$$Ax = b, \quad x \geq 0$$

a basic solution partitions the variables into two sets:

$$\begin{bmatrix} B & N \end{bmatrix} \begin{bmatrix} x_B \\ x_N \end{bmatrix} = b \quad \text{with} \quad x_N = 0$$

ie,

$$Bx_B = b$$

a basic feasible solution in standard form is **degenerate** if more than $n - m$ components in x are zero, ie,

$$x = \begin{bmatrix} x_B \\ x_N \end{bmatrix} \begin{matrix} m \\ n - m \end{matrix} = \begin{bmatrix} B^{-1}b \\ 0 \end{bmatrix} \leftarrow \text{has some zero components}$$