

Proofs of Conditional Statements

- Many theorems have the form:
 - *For all x , $H(x) \rightarrow C(x)$*
- The cube of every negative real number is negative.
 - For all x , $(x^3 < 0) \rightarrow (x < 0)$
 - Domain: set of real numbers.
- If x^2 is an odd integer, then x is also an odd integer.
 - For all x , $(x^2 \text{ is odd}) \rightarrow (x \text{ is odd})$
 - Domain: set of integers.

Ways to prove conditional statements

- *Prove*: $\forall x H(x) \rightarrow C(x)$
- Direct proof
 - Name a generic element x in the domain that satisfies $H(x)$
 - Prove that $C(x)$ is also true
- Implicitly uses universal instantiation.

Ways to prove conditional statements

- *Prove:* $\forall x H(x) \rightarrow C(x)$
- Proof by contrapositive
 - Name a generic element x in the domain such that $C(x)$ is false
 - Prove that $H(x)$ is also false
 - At the beginning of the proof, state what you are assuming ($C(x)$ is false) and state what you will prove ($H(x)$ is false).
- *Proofs by contrapositive work because:*
 - $\forall x H(x) \rightarrow C(x)$ is logically equivalent to
$$\forall x \neg C(x) \rightarrow \neg H(x)$$

Direct proof example

- Theorem: The cube of every even integer is even.
- Proof:
 - Let x be an even integer.
 - We will show that x^3 is also even.
 - Then $x = 2k$ for some integer k .
 - Cube both sides of the equation to get:
 - $x^3 = (2k)^3 = 2^3k^3 = 8k^3 = 2(4k^3)$
 - $4k^3$ is an integer, so x^3 can be expressed as 2 times an integer, and therefore x^3 is even.

Direct proof example

- Theorem: The product of two rational numbers is rational.
- Proof:
 - Let x and y be rational numbers.
 - We will show that xy is also rational.
 - Then $x = a/b$ and $y = c/d$, where a, b, c, d are integers and $b \neq 0$ and $d \neq 0$.
 - Multiply x and y to get:
 - $xy = \left(\frac{a}{b}\right) \left(\frac{c}{d}\right) = \left(\frac{ac}{bd}\right)$
 - ac and bd are both integers. Also $bd \neq 0$ because $b \neq 0$ and $d \neq 0$.
 - Therefore xy is the ratio of two integers with a non-zero denominator which means that xy is rational.

Proof by contrapositive example.

- Theorem: If n is an integer and $3n+7$ is odd then n is even.
- Proof:
 - Let n be an arbitrary integer such that n is odd.
 - We will show that $3n+7$ is even.
 - $n = 2k+1$, for some integer k .
 - Plug in the expression $n=2k+1$ into $3n+7$
 - $3n+7 = 3(2k+1)+7 = 6k+3+7 = 6k+10 = 2(3k+5)$
 - Since k is an integer, then $3k+5$ is also an integer. Therefore $3n+7$ can be expressed as 2 times an integer and therefore $3n+7$ is even.

Proof by contrapositive example.

- Theorem: If x is a real number such that $3x$ irrational then x is irrational.
- Before we start:
 - Every real number is rational or irrational.
 - A real number is rational if it can not be expressed as the ratio of two integers.

Proof by contrapositive example.

- Theorem: If x is a real number such that $3x$ irrational then x is irrational.
- Proof
 - Assume that x is a real number and that x is not irrational.
 - Will show that $3x$ is rational and therefore not irrational.
 - Since x is real and not irrational, then it is rational.
 - $x = \frac{a}{b}$, where $b \neq 0$.
 - Therefore $3x = 3 \frac{a}{b} = \frac{3a}{b}$.
 - Since a is an integer $3a$, is also an integer. Furthermore $b \neq 0$. Therefore $3x$ can be expressed as the ratio of two integers with a non-zero denominator which means that $3x$ is rational.

Idempotent laws:

$$p \vee p \equiv p$$

$$p \wedge p \equiv p$$

Associative laws:

$$(p \vee q) \vee r \equiv p \vee (q \vee r)$$

$$(p \wedge q) \wedge r \equiv p \wedge (q \wedge r)$$

Commutative laws:

$$p \vee q \equiv q \vee p$$

$$p \wedge q \equiv q \wedge p$$

Distributive laws:

$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$

Identity laws:

$$p \vee F \equiv p$$

$$p \vee T \equiv T$$

$$p \wedge F \equiv F$$

$$p \wedge T \equiv p$$

Involution law:

$$\neg\neg p \equiv p$$

Complement laws:

$$p \vee \neg p \equiv T$$

$$\neg T \equiv F$$

$$p \wedge \neg p \equiv F$$

$$\neg F \equiv T$$

De Morgan's laws:

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

Conditional identities:

$$p \rightarrow q \equiv \neg p \vee q$$

$$p \leftrightarrow q \equiv (p \rightarrow q) \wedge (q \rightarrow p)$$